

# On $K$ -ary $n$ -Cubes ad Isometric Words<sup>\*</sup>

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A  $k$ -ary  $n$ -cube is a graph with  $k^n$  vertices, each associated to a word of length  $n$  over an alphabet of cardinality  $k$ . The subgraph obtained deleting those vertices which contain a given  $k$ -ary word  $f$  as a factor is called the  $k$ -ary  $n$ -cube avoiding  $f$ . When, for any  $n$ , such a subgraph is *isometric* to the cube, the word  $f$  is said *isometric* (or *good*). In the binary case, isometric words can be equivalently defined, independently from hypercubes. Recently, this problem has been investigated in the binary case [4]. These two approaches are here considered and extended to the  $k$ -ary alphabet, showing that they are still coincident for  $k = 3$ , also for  $k = 4$ , not with Hamming distance but with Lee distance, and they are not still coincident from  $k > 4$  on. [1, 2]. A word  $f$  is also called  $d$ -good if for any pair of words  $u$  and  $v$  of length  $d$ , with  $d \geq |f|$ , which do not contain the factor  $f$ ,  $u$  can be transformed in  $v$  by exchanging one by one the symbols on which they differ and generating only words which do not contain  $f$ . It is *good* if it is  $d$ -good for all  $d \geq |f|$ , it is *bad* if it is not good. The *index* of a bad word is the threshold  $d$  from which the word is no longer  $d$ -good, and a pair of words  $(u, v)$  showing that the word is not good is called a pair of *witnesses* for the bad word. A characterization of these  $k$ -ary words is here provided also in terms of *2-error-overlap*. Furthermore, here is an algorithm to decide whether a word is bad or not. When the word is bad the algorithm provides its index and a pair of witnesses of minimal length. An analogous algorithm is given for binary alphabet in [4]; it runs in cubic time. Our algorithm works on  $k$ -ary words, with  $k \leq 4$  and construct an enhanced suffix tree of the word. This data structure has been used in [3] to design an algorithm to decide whether a  $k$ -ary word is “bad” or not (without computing the index and the witnesses). Algorithm here provided runs in  $O(n)$  time with a preprocessing of  $O(n)$  time and space.

## References

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